

# Discrete Model of TCP Congestion Control Algorithm with Round Dependent Loss Rate

O. I. Bogoiavlenskaia  
**Presenter** Dmitry Korzun

*Petrozavodsk State University*

*olbgvl@cs.karelia.ru*

# Motivation and Problem Statement

- Internet needs a tool to control its performance and resource sharing
- This service is provided at the end-to-end level by Transmission Control Protocol (TCP)
- TCP performs distributed flow control. It controls **performance** of the network to prevent it from **congestion** collapse.

Key **metrics** of congestion control performance:

- Average congestion window size
- Variance of of congestion window size

# Motivation and Problem Statement

Application areas:

- Performance Evaluation and Capacity Planning of networking links and networks
- Throughput evaluation across end-to-end path
- Performance prediction for jitters sensitive applications (variance), e.g. media streams
- Choose cloud provider for fetching contents
- Support network administration and design
- Provides foundation for further research and development on transport layer protocols.

# Distributed Flow Control

- Sliding window is amount of data which sender is allowed to inject in the network without acknowledgment
- Flow control means control on sliding window size. TCP uses set of algorithms to control its window size  $W$

Additive Increase Multiplicative Decrease Algorithm (AIMD)

$$\begin{array}{ccc} W & \xrightarrow{\text{delivery}} & W + 1 \\ \downarrow \text{loss} & & \\ W\alpha & & \end{array}$$

$\alpha < 1$  is rational number  $\alpha = \frac{n}{m}$ ,  $n < m$

# Details on Sliding Window Size

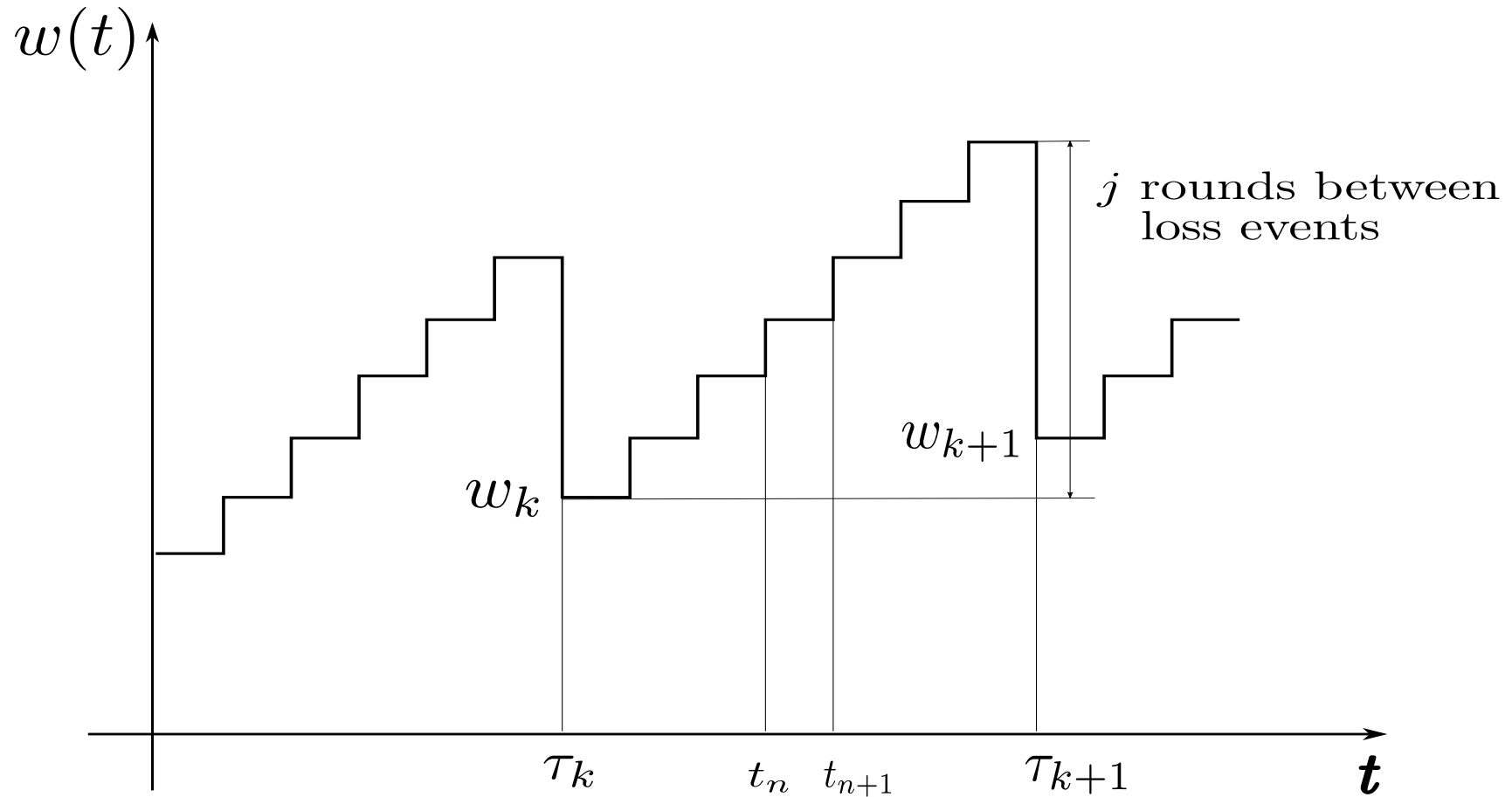


Figure 1: Evolution of TCP congestion window size

## Main Definitions

Let  $w(t)$  be cwnd size,  $t_n$  are AIMD-rounds end-points.

Let  $\tau_k$  be equal to the first moment  $t_n$ , arrived after  $k$ th data loss event, i.e.

$$\tau_k = t_n : w(t_n + 0) = \left\lfloor \frac{w(t_n)n}{m} \right\rfloor, \quad k = 1, 2, \dots$$

and  $\tau_{k_1} < \tau_{k_2}$ , if  $k_1 < k_2$ .

Discrete distribution  $\{f_j\}_{j \geq 0}$ , i. e. the distribution of the number of the consequential rounds at the end of which  $w(t)$  size increased.

In other words  $f_j$  is the probability that a sequence of  $j$  rounds without losses took place after a segment loss event occurred.

## Ergodic Properties

According to the assumptions and definitions above the sequence  $\{w_k = w(\tau_k + 0)\}_{k>0}$  is the **Markov chain**.

Let's define the expectation determined by  $f_j$  as  $R = \sum_{j=1}^{\infty} j f_j$ .

**Theorem 1** *If  $R$  is finite then the Markov chain  $\{w_k\}$  has steady state distribution.*

## The Expectation Estimate

$$w_{k+1} = \left\lfloor \frac{(w_k + r_k)n}{m} \right\rfloor, \quad (1)$$

where  $r_k$  is the random variable with the distribution  $\{f_j\}$ .

$$w_{k+1} = \frac{(w_k + r_k)n - i_k}{m}, \quad (2)$$

where  $i_k$  is the integer random variable and  $0 \leq i_k < m$ . Difference equation (2) has a stationary solution

$$w_k^* = \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^{j+1} r_{k-j} - \frac{1}{m} \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^j i_{k-j}. \quad (3)$$



## The Expectation Estimate

**Theorem 2** *If the inequality  $R > \frac{m}{n}$  holds then the stationary expectation of the congestion window size can be estimated as follows*

$$\mathbf{E}[w_k^*] \geq R \frac{n}{m-n} - \frac{m-1}{m-n}. \quad (4)$$

P R O O F. Processing of the expression (3) yields the following

$$\mathbf{E}[w_k^*] = \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^{j+1} \mathbf{E}[r_{k-j}] - \frac{1}{m} \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^j \mathbf{E}[i_{k-j}]. \quad (5)$$

By definitions  $\forall k, j$   $\mathbf{E}[r_{k-j}] = R$  and  $\mathbf{E}[i_{k-j}] \leq m-1$ . Hence

$$\mathbf{E}[w_k^*] \geq R \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^{j+1} - \frac{m-1}{m} \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^j = R \frac{n}{m-n} - \frac{m-1}{m-n}.$$

The last expression is positive if  $Rn > m$ .  $\square$

## Second Moment Estimate

$$\mathbb{E} \left[ w_k^{*2} \right] = \mathbb{E} \left[ \left( \sum_{j=0}^{\infty} \binom{n}{m}^{j+1} r_{k-j} \right)^2 \right] - \quad (W_1)$$

$$-\frac{2}{m} \mathbb{E} \left[ \sum_{j=0}^{\infty} \binom{n}{m}^{j+1} r_{k-j} \sum_{s=0}^{\infty} \binom{n}{m}^s i_{k-s} \right] + \quad (W_2)$$

$$+\frac{1}{m^2} \mathbb{E} \left[ \left( \sum_{j=0}^{\infty} \binom{n}{m}^j i_{k-j} \right)^2 \right] \quad (W_3)$$

## Second Moment Estimate

Let's  $\forall k, j$   $R^{(2)} = \mathbf{E}[r_{k-j}^2]$  and  $\forall k, j \neq s$   $R^2 = \mathbf{E}[r_{k-j}r_{k-s}]$ . Then

$$\begin{aligned}
 W_1 &= R^{(2)} \sum_{j=1}^{\infty} \left(\frac{n}{m}\right)^{2j} + 2R^2 \sum_{j=0}^{\infty} \sum_{s=j+1}^{\infty} \left(\frac{n}{m}\right)^{j+s+2} = \\
 &R^{(2)} \frac{n^2}{m^2 - n^2} + 2R^2 \frac{n^3}{(m-n)(m^2 - n^2)} \tag{6}
 \end{aligned}$$

## Second Moment Estimate

$$\begin{aligned}
 W_2 &= \frac{2}{m} \sum_{j=0}^{\infty} \sum_{s=0}^{\infty} \left(\frac{n}{m}\right)^{j+s+1} \mathbf{E}[r_{k-j} i_{k-s}] < \\
 2R \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^{j+1} \left(1 - \frac{n}{m}\right)^{-1} &= \frac{2Rmn}{(m-n)^2} = E_l. \tag{7}
 \end{aligned}$$

Furthermore  $\forall k, j \mathbf{E}[i_{k-j}^2] < m^2$ . Also  $\forall k, j \neq s \mathbf{E}[i_{k-j} i_{k-s}] < m^2$  as well.

$$\begin{aligned}
 W_3 &< \sum_{j=0}^{\infty} \left(\frac{n}{m}\right)^{2j} + 2 \sum_{j=0}^{\infty} \sum_{s=j+1}^{\infty} \left(\frac{n}{m}\right)^{j+s} = \\
 \frac{m^2}{m^2 - n^2} + \frac{2nm^2}{(m-n)(m^2 - n^2)} &= E_g. \tag{8}
 \end{aligned}$$

## Second Moment Estimate

**Theorem 3** *If*

$$\frac{n^2 R}{m+n} > m. \quad (9)$$

*then the following interval estimate holds*

$$W_1 - E_l < \mathbf{E}[w_k^{*2}] < W_1 + E_g. \quad (10)$$

### Compare with

Altman E., Avrachenkov K., Barakat C. A Stochastic model of TCP/IP with Stationary Random Losses, Proceedings of ACM SIGCOMM'00. Stockholm, 2000. pp. 231-242.

The term  $W_1$  coincides up to notation and the terms  $E_l$  and  $E_g$  appear due to the discrete nature of the random process considered and due to the using of the *floor* operation in the AIMD control.

## Conclusion

- The stepwise model of AIMD New Reno congestion control with **arbitrary** decrease coefficient is analyzed.
- The theorem on its ergodic properties is proved.
- The estimate of stationary **expectation** of congestion window size is obtained and its applicability terms are derived.
- The interval estimates of stationary congestion window size **second moment** is obtained and its applicability terms are derived.
- Further analysis yields bounds of the standard deviation which could not be obtained through expectation estimates constructed using Goelder's inequality.