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What is Nonnegative Linear Diophantine Equation (NLDE)?	Hilbert basis of NLDE system
<b>What is Nonnegative Linear Diophantine Equation (NLDE)?</b>	$\frac{\text{Hilbert basis of NLDE system}}{\text{Consider NLDE system: } Ax = b. \text{ There exist finite sets } \mathcal{N} \text{ and } \mathcal{H}:}$
<i>n</i> equations, <i>m</i> unknowns: $Ax = b$ $A \in \mathbb{Z}^{n \times m}, \ b \in \mathbb{Z}^n, \ x \in \mathbb{Z}^m$	Consider NLDE system: $Ax = b$ . There exist finite sets $\mathcal{N}$ and $\mathcal{H}$ : $\mathcal{N} = \{x^{(1)}, x^{(2)}, \dots, x^{(p)}\},  \mathcal{H} = \{h^{(1)}, h^{(2)}, \dots, h^{(q)}\}$
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$n \text{ equations, } m \text{ unknowns: } Ax = b  A \in \mathbb{Z}^{n \times m}, \ b \in \mathbb{Z}^n, \ x \in \mathbb{Z}^m$ $\begin{cases} a_{11}x_1 + \dots + a_{1m}x_m = b_1 \\ \dots \\ a_{n1}x_1 + \dots + a_{nm}x_m = b_n \end{cases}$ $\text{Diophantness: solutions are integer, } x_i \in \{0, 1, 2, \dots\} = \mathbb{Z}_+$	Consider NLDE system: $Ax = b$ . There exist finite sets $\mathcal{N}$ and $\mathcal{H}$ : $\mathcal{N} = \{x^{(1)}, x^{(2)}, \dots, x^{(p)}\},  \mathcal{H} = \{h^{(1)}, h^{(2)}, \dots, h^{(q)}\}$ $\boxed{\mathcal{S} = \mathcal{N} + \mathcal{H}^*}$ $x \in \mathcal{S} \iff x = x^{(l)} + \alpha_1 h^{(1)} + \dots + \alpha_q h^{(q)} = x^{(0)} + \sum_{s=1}^q \alpha_s h^{(s)}$ for some $x^{(l)} \in \mathcal{N}$ and $\alpha_s \in \mathbb{Z}_+$ Solution $x$ is <i>minimal</i> if there is no solution $x'$ such that $x' \leq x$
$n \text{ equations, } m \text{ unknowns: } Ax = b  A \in \mathbb{Z}^{n \times m}, \ b \in \mathbb{Z}^n, \ x \in \mathbb{Z}^m \\ \begin{cases} a_{11}x_1 + \dots + a_{1m}x_m = b_1 \\ \dots \\ a_{n1}x_1 + \dots + a_{nm}x_m = b_n \end{cases}$ $Diophantness:  solutions are integer, \ x_i \in \{0, 1, 2, \dots\} = \mathbb{Z}_+$ $Nonnegativness:  solutions are nonnegative, \ x_i \ge 0$ $Example:$	Consider NLDE system: $Ax = b$ . There exist finite sets $\mathcal{N}$ and $\mathcal{H}$ : $\mathcal{N} = \{x^{(1)}, x^{(2)}, \dots, x^{(p)}\},  \mathcal{H} = \{h^{(1)}, h^{(2)}, \dots, h^{(q)}\}$ $\boxed{\mathcal{S} = \mathcal{N} + \mathcal{H}^*}$ $x \in \mathcal{S} \iff x = x^{(l)} + \alpha_1 h^{(1)} + \dots + \alpha_q h^{(q)} = x^{(0)} + \sum_{s=1}^q \alpha_s h^{(s)}$ for some $x^{(l)} \in \mathcal{N}$ and $\alpha_s \in \mathbb{Z}_+$
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$n \text{ equations, } m \text{ unknowns:} \qquad Ax = b \qquad A \in \mathbb{Z}^{n \times m}, \ b \in \mathbb{Z}^n, \ x \in \mathbb{Z}^m \\ \begin{cases} a_{11}x_1 + \dots + a_{1m}x_m = b_1 \\ \dots \\ a_{n1}x_1 + \dots + a_{nm}x_m = b_n \end{cases}$ $Diophantness:  \text{solutions are integer, } x_i \in \{0, 1, 2, \dots\} = \mathbb{Z}_+$ $Nonnegativness:  \text{solutions are nonnegative, } x_i \ge 0$ $Example: \qquad \begin{cases} x_1 - x_2 + 3x_3 = 1 \\ x_1 + 2x_2 - x_4 = -1 \end{cases} \qquad \begin{cases} x_1 - x_2 + 3x_3 = 0 \\ x_1 + 2x_2 - x_4 = -1 \end{cases} \qquad \begin{cases} x_1 - x_2 + 3x_3 = 0 \\ x_1 + 2x_2 - x_4 = 0 \end{cases}$	Consider NLDE system: $Ax = b$ . There exist finite sets $\mathcal{N}$ and $\mathcal{H}$ : $\mathcal{N} = \{x^{(1)}, x^{(2)}, \dots, x^{(p)}\},  \mathcal{H} = \{h^{(1)}, h^{(2)}, \dots, h^{(q)}\}$ $\boxed{\mathcal{S} = \mathcal{N} + \mathcal{H}^*}$ $x \in \mathcal{S} \iff x = x^{(l)} + \alpha_1 h^{(1)} + \dots + \alpha_q h^{(q)} = x^{(0)} + \sum_{s=1}^q \alpha_s h^{(s)}$ for some $x^{(l)} \in \mathcal{N}$ and $\alpha_s \in \mathbb{Z}_+$ Solution $x$ is minimal if there is no solution $x'$ such that $x' \leq x$ (component-wise partial order: $x'_i \leq x_i, i = 1, 2, \dots, m$ ) $\mathcal{N}$ all minimal solutions of $Ax = b$ , $\mathcal{H}$ all minimal solutions of $Ax = 0$ (homogenous system),
$n \text{ equations, } m \text{ unknowns: } Ax = b  A \in \mathbb{Z}^{n \times m}, \ b \in \mathbb{Z}^n, \ x \in \mathbb{Z}^m \\ \begin{cases} a_{11}x_1 + \dots + a_{1m}x_m = b_1 \\ \dots \\ a_{n1}x_1 + \dots + a_{nm}x_m = b_n \end{cases}$ $Diophantness:  solutions are integer, \ x_i \in \{0, 1, 2, \dots\} = \mathbb{Z}_+$ $Nonnegativness:  solutions are nonnegative, \ x_i \ge 0$ $Example:$	Consider NLDE system: $Ax = b$ . There exist finite sets $\mathcal{N}$ and $\mathcal{H}$ : $\mathcal{N} = \{x^{(1)}, x^{(2)}, \dots, x^{(p)}\},  \mathcal{H} = \{h^{(1)}, h^{(2)}, \dots, h^{(q)}\}$ $\boxed{\mathcal{S} = \mathcal{N} + \mathcal{H}^*}$ $x \in \mathcal{S} \iff x = x^{(l)} + \alpha_1 h^{(1)} + \dots + \alpha_q h^{(q)} = x^{(0)} + \sum_{s=1}^q \alpha_s h^{(s)}$ for some $x^{(l)} \in \mathcal{N}$ and $\alpha_s \in \mathbb{Z}_+$ Solution $x$ is <i>minimal</i> if there is no solution $x'$ such that $x' \leq x$ (component-wise partial order: $x'_i \leq x_i, i = 1, 2, \dots, m$ ) $\mathcal{N}$ all minimal solutions of $Ax = b$ ,

NLDE system of 2 equations in 4 unknowns and its homogenous case:

$$\begin{cases} x_1 - x_2 + 3x_3 = 1\\ x_1 + 2x_2 - x_4 = -1 \end{cases} \qquad \begin{cases} x_1 - x_2 + 3x_3 = 0\\ x_1 + 2x_2 - x_4 = 0 \end{cases}$$
$$\mathcal{N} = \begin{cases} \begin{pmatrix} 1\\0\\0\\2 \end{pmatrix}, \begin{pmatrix} 0\\2\\1\\5 \end{pmatrix} \end{cases}, \quad \mathcal{H} = \begin{cases} \begin{pmatrix} 1\\1\\0\\3 \end{pmatrix}, \begin{pmatrix} 0\\3\\1\\6 \end{pmatrix} \end{cases}$$
$$\begin{cases} 2\\1\\6 \end{pmatrix} \end{cases}$$
$$\begin{pmatrix} 2\\1\\0\\2 \end{pmatrix} + \begin{pmatrix} 1\\4\\1\\9\\2 \end{pmatrix} = \begin{pmatrix} 1\\0\\0\\2 \end{pmatrix} + 1\begin{pmatrix} 1\\1\\0\\3\\1 \end{pmatrix} + 1\begin{pmatrix} 0\\3\\1\\6 \end{pmatrix} = \begin{pmatrix} 0\\2\\1\\5 \end{pmatrix} + 2\begin{pmatrix} 1\\1\\0\\3 \end{pmatrix}$$

Solvability: has NLDE a solution?

Particular solution: find a solution

#### Hilbert basis:

- search the basis
- how many elements in the basis?
- given solution = a basis solution ?
- given set of solutions = Hilbert basis?

• . . .

# **Complexity of Solving NLDE**

The source of a rich family of complexity problems

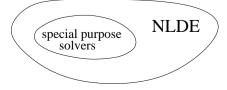
Solvability & Particular solution: has NLDE a solution? find a solution

- Polynomial for homogenous NLDE
- NP-complete in general case

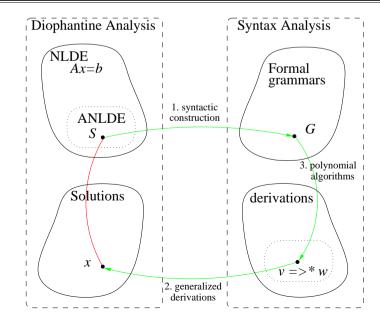
#### Hilbert basis: over-NP in general case

- search the basis  $-|\mathcal{H}|$  depends exponentially on size (n, m, ||A, b||)
- how many elements in the basis? #P-hard and in #NP
- given solution = a basis solution? coNP-complete
- given set of solutions = Hilbert basis? is the set is basis of some NLDE system?

Universal solvers are not adequate for practical use!



# Interrelationship: NLDE $\leftrightarrow$ Formal grammars



Associated with formal grammar ANLDE system		Associated with formal grammar ANLDE system
Grammar G and strings $v, w$ $m$ rules $(r_i)$ $n$ nonterminals $(A_k)$ $t$ terminals $(a_k)$ ANLDE system $Ax = b$ $m$ unknowns $(x_i)$ n + t equations $(k)x$ — the number of applications of gramm A — rules structure (occurrences of symbol b — $v, w$ structure (occurrences of symbol	ols)	$\begin{split} \sum_{i=1}^{m} \operatorname{occ}(U, \operatorname{lhs}(r_i)) x_i &= \sum_{i=1}^{m} \operatorname{occ}(U, \operatorname{rhs}(r_i)) x_i = \operatorname{occ}(U, v) - \operatorname{occ}(U, w) ,  U \in N \cup T \\ & \operatorname{occ}(U, \alpha)  \text{number of occurrences of symbol } U \text{ in string } \alpha \\ & \operatorname{lhs}(r) \text{ or rhs}(r) \text{ left or right hand side of rule } r \\ & r_1 \colon A \to AAB  m = 4, n = 2, t = 0 \\ & r_2 \colon A \to BB  v = B, w = A \\ & r_3 \colon B \to AAAB  B \Rightarrow^* A? \\ & r_1 \colon B \to \varepsilon  (\text{hom: } A \Rightarrow^+ A \text{ and } B \Rightarrow^+ B) \\ & A \colon (x_1 + x_2) - (2x_1 + 3x_3) = -1 \\ & B \colon (x_3 + x_4) - (x_1 + x_2 + x_3) = 1 \\ & x = (0, 2, 1, 5) \colon B \stackrel{3}{\Rightarrow} AAAB \stackrel{4}{\Rightarrow} AAA \stackrel{22}{\Rightarrow} ABBBB \stackrel{444}{\Rightarrow} A \\ & x = (1, 0, 0, 2) \colon BA \stackrel{1}{\Rightarrow} BAAB \stackrel{44}{\Rightarrow} AAB \\ & h = (1, 1, 0, 3) \colon A \stackrel{1}{\Rightarrow} AAB \stackrel{2}{\Rightarrow} ABBB \stackrel{444}{\Rightarrow} A \end{split}$
standard derivation $v \Rightarrow^* w$ $\rightarrow$ (b)	$\begin{array}{c} \text{nd Derivations} \\ \text{by construction }) \rightarrow \\ \leftarrow (?) \leftarrow \end{array}  \text{solution } x \in \mathbb{Z}_{+}^{m} \\ \\ \Rightarrow^{*} \widetilde{w} \longleftrightarrow  \text{solution } x \in \mathbb{Z}_{+}^{m} \end{array}$	Homogenous ANLDE Systems $\sum_{i \in H_k} x_i = \sum_{i=1}^n \gamma_{ki} x_i$ , $k = 1, 2,, n$ , $I_1 \cup \cup I_n = \{1, 2,, m\}$ Polynomial complexity $(n \le m)$ :

Three-component solution:

$$x = x^{v,w'} + x^{w \setminus w'} + x^{\varepsilon}$$

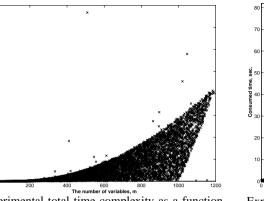
Generalized derivation:

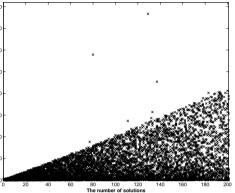
Hilbert basisParticular solutionTime
$$O(q^3m^2n^2) = O(q^3m^4)$$
 $O(m^2n^2) = O(m^4)$ Space $O(qmn) = O(qm^2)$  $O(mn^2) = O(m^3)$ 

**Experimentally:** time is  $\Theta(qm^2)$  in the most test cases

- more than 1 billion of test ANLDE systems
- coefficients  $\gamma_{ki}$ : ..10<sup>5</sup>
- unknowns and equations:  $..10^3$

## **Experimental plots**





Experimental total time complexity as a function on the number of unknowns

Experimental total time complexity as a function on the number of solutions

The dimensions of the generated systems were in range:  $n \in [1, 1000]$ ,  $m \in [n, n + 200]$ ,  $\gamma_{ki} \in [0, 500]$ . The systems were generated in such a way that each of them has at least one basis solution but not more than 200.

### Experimental analysis and testing

- 1. Think up an ANLDE system and solve it estimating the resources (time&space), test the found solution
- 2. Construct a generator, generate a system and its solution, ..., compare the found and original solutions
- 3. Compare the results with an alternative solver
- 4. More detailed analysis needs not a system, but a set of them: generate a set and analyze how the solver solves all the systems

Combine these features together, move this to the web space . . .

#### AND YOU GET THE WEB-SYNDIC SOFTWARE SYSTEM

Conclusion

THE WEB-SYNDIC SYSTEM IS A TOPIC TO TALK FOR THE NEXT PRESENTATION . . .

# Thank you! Any questions?