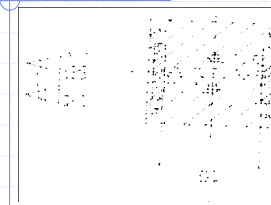


# Efficient algorithms for polygonal approximation

Alexander Kolesnikov

## Examples of polygonal approximation



Vectorization

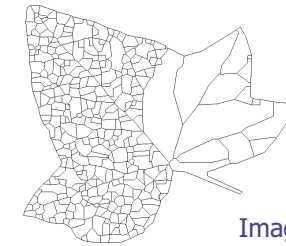
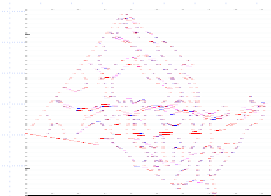
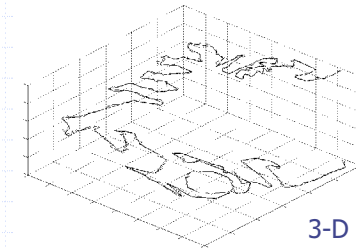


Image analysis



Digital cartography



3-D paths

## Heuristic algorithms for approximation

- 1) Sequential tracing approach
- 2) Split method
- 3) Merge method
- 4) Split-and-Merge method
- 5) Dominant point detection
- 6) Relaxation labeling
- 7)  $K$ -means method
- 8) Genetic (evolutional) algorithms
- 9) Ant colony optimization method
- 10) Tabu search
- 11) Discrete particle swarm algorithm
- 12) Vertex adjustment method

## Min- $\epsilon$ problem: Motivation

### Heuristic algorithms

Non-optimal:  $F < 100\%$

Fast:  $O(N) - O(N^2)$   
(seconds and less)

### Optimal algorithm

Optimal:  $F = 100\%$

Slow:  $O(N^2) - O(N^3)$   
(minutes and more)

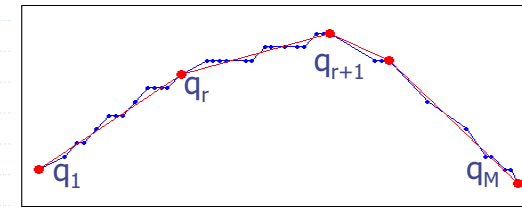
## Min- $\epsilon$ problem: Motivation

|  |  |   |
|--|--|---|
| <b>Heuristic algorithms</b><br>Non-optimal: $F < 100\%$<br>Fast: $O(N) - O(N^2)$<br>(seconds and less) |  | <b>Optimal algorithm</b><br>Optimal: $F = 100\%$<br>Slow: $O(N^2) - O(N^3)$<br>(minutes and more) |
|--|--|---|

### Efficient algorithm

Fast:  $O(N) - O(N^2)$   
 Close to optimal:  $F \approx 100\%$

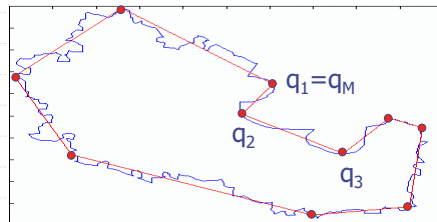
## Min- $\epsilon$ problem for open curve



Approximate the given open  $N$ -vertex polygonal curve  $P$  by another one  $Q$  consisting of at most  $M$  line segments with minimum error  $E(P)$ :

$$E(P) = \min_{\{q_r\}} \sum_{r=2}^{M-2} e^2(q_r, q_{r+1})$$

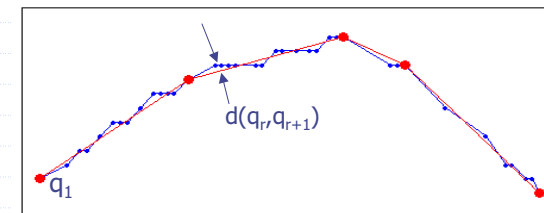
## Min- $\epsilon$ problem for closed curves



Approximate the given closed  $N$ -vertex polygonal curve  $P$  by another one  $Q$  consisting of at most  $M$  line segments with minimum error  $E(P)$ :

$$E(P) = \min_{\{q_r\}} \sum_{r=1}^{M-1} e^2(q_r, q_{r+1})$$

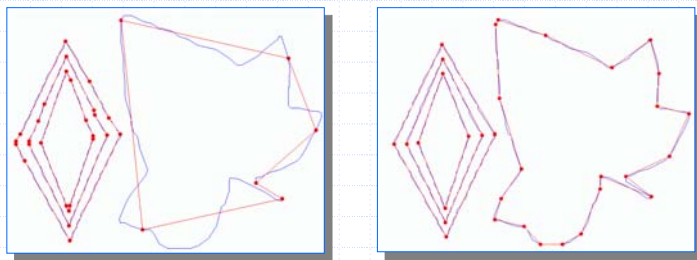
## Min-# problem for open curve



Given polygonal curve  $P$ , approximate it by another polygonal curve  $Q$  with the minimum number of segments  $M$  so that the approximation error (distortion) does not exceed a given maximum tolerance  $\epsilon$ :

$$M \rightarrow \min \quad \text{subject to: } d(P) \leq \epsilon$$

## Multi-object $\min\text{-}\varepsilon$ problem



Given  $K$  polygonal curves  $P_1, P_2, \dots, P_K$ , approximate it by set of  $K$  another polygonal curves  $Q_1, Q_2, \dots, Q_K$  with a given total number of segments  $M$ .

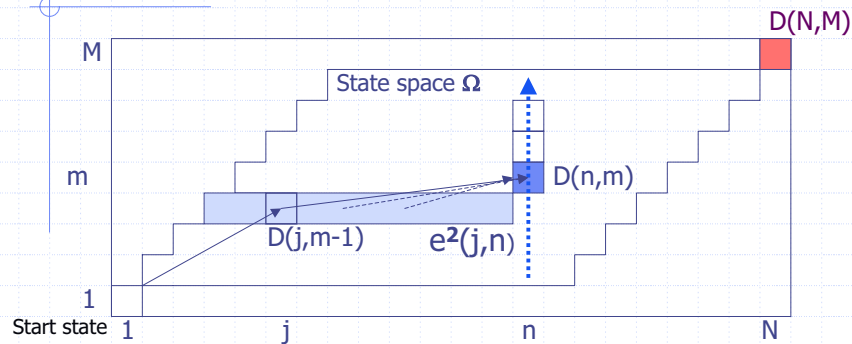
## Multi-object $\min\text{-}\varepsilon$ problem (cont'd)

Given  $K$  polygonal curves  $P_1, P_2, \dots, P_K$ , approximate it by set of  $K$  another polygonal curves  $Q_1, Q_2, \dots, Q_K$  with a given total number of segments  $M$  so that the total approximation error with measure  $L_2$  is minimized.

$$E(P_1, \dots, P_K, M) = \min_{\{M_k\}} \min_{\{q_m\}} \sum_{k=1}^K \sum_{m=1}^{M_k-1} e^2(q_{k,m}, q_{k,m+1})$$

subject to:  $\sum_{k=1}^K M_k \leq M$

## DP algorithm for $\min\text{-}\varepsilon$ problem



Dynamic programming algorithm:

$$D(n, m) = \min \{ D(j, m-1) + e^2(j, n) \mid m \leq j < n \}$$

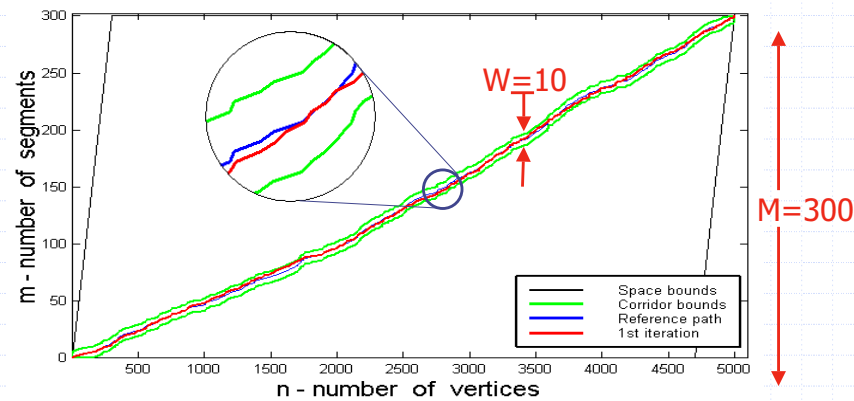
$$\forall (n, m) \in \Omega$$

## Iterative reduced search for $\min\text{-}\varepsilon$ problem

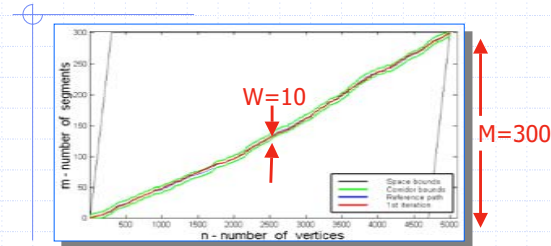
- Step 1: Find preliminary approximation with any fast heuristic algorithm.
- Step 2: Construct bounding corridor in state space along the reference path
- Step 3: Perform DP search in the bounding corridor

Use the output as a reference solution and repeat the search

## Iterative reduced search DP algorithm



## Iterative reduced search DP algorithm



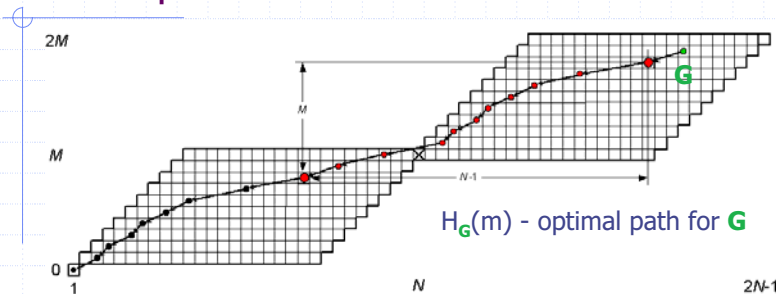
Complexity:  $O(W^2N^2/M) = O(N) - O(N^2)$

Speed-up:  $(W/M)^2$

Example:  $(W/M)^2 = (10/300)^2 = 1/900$

Fidelity: up to 100% for iterative search

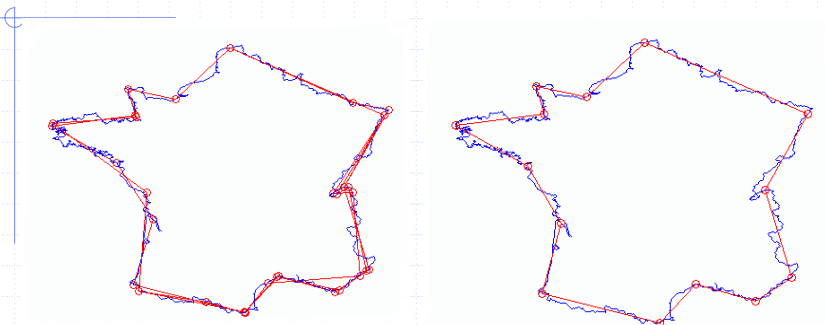
## Min- $\varepsilon$ problem for closed curves



Conjugate states:  $H_G(m-M) = H_G(m) - (N-1)$ ;

$n_{\text{start}} = \arg \min_{M \leq m \leq 2M} \{D(H_G(m), m) - D(H_G(m-M), m-M)\} - (N-1)$ ;

## Min- $\varepsilon$ problem for closed curves: results



Heuristic algorithm:  
Fidelity  $F=88-100\%$   
 $T=100$  s

Proposed algorithm:  
Fidelity  $F=100\%$   
 $T=10.4$  s.

## Iterative reduced search for multi-object min- $\varepsilon$ problem

**Step 1:** Find preliminary approximation of every object for an initial number of segments.

**Step 2:** Iterate the following:

- Apply multiple-goal reduced search DP to define the Rate-Distortion functions.
- Solve the optimal allocation of the segments number among the objects using the Rate-Distortion functions.

## Min-# solution as shortest path in digraph

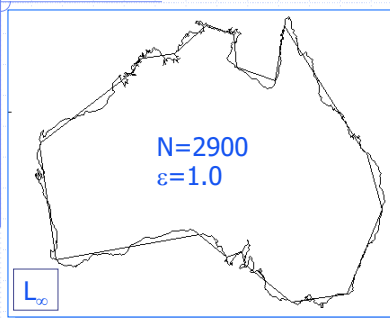


Feasibility graph  $G_\varepsilon(P)$   
for  $\varepsilon=10$  with  $L_\infty$  measure.



Min-# solution as shortest  
path in the digraph  $G_\varepsilon(P)$ .

## Joint using of $L_\infty$ & $L_2$ error measures



A shortest path in digraph  $G_\varepsilon(P)$   
for  $\varepsilon=1.0$ :  $E_2=629$ ,  $d_{\max}=1.0$ ;  
 $T=4.5$  s.



After optimization with reduced  
search:  $E_2=298$ ,  $d_{\max}=1.08$ ;  
 $T=5.0$  s.

## Main results

- Iterative Reduced Search algorithm, including:
  - Algorithm for min- $\varepsilon$  problem
  - Algorithm for multiple-object min- $\varepsilon$  problem
  - Algorithm for min-# problem with  $L_\infty$  &  $L_2$  error measures
- Approximation algorithm for closed curves with state space analysis

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URL: [http://cs.joensuu.fi/~koles/approximation/Ch3\\_0.html](http://cs.joensuu.fi/~koles/approximation/Ch3_0.html)